Statically Bounded-Memory Delayed Sampling for Probabilistic Streams



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- **Goal:** Run inference in *bounded memory*.
- I.e. even if programs run for infinite time, should have a finite memory footprint

Bounded-Memory Delayed Sampling^{1,2}



- **Goal:** Bound the size of the delayed sampling graph.
- Maintain a finite set of *reachable* nodes that are accessible from pointers in the program state

1: Murray et. al. "Delayed Sampling and Automatic Rao—Blackwellization of Probabilistic Programs". AISTATS 2018 2: Baudart et. al. "Reactive Probabilistic Programming". PLDI 2020

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Semantic Properties

- We can define properties on *traces* of probabilistic programs
- A trace records all operations the program executes

let kalman = stream {	xЮ
init = 0.0;	ye
<pre>step (pre x, obs) = {</pre>	ob
let $x = sample(gaussian(pre x.1))$ in	x1
$\int \frac{d}{dt} = \int \frac$	y1
<pre>let y = sample(gaussian(x,1)) in</pre>	ob
observe (y, obs);	x2
(\mathbf{x}, \mathbf{x})	v2
	y Z
۲ ۲	OD
}	• •

The *m*-consumed Property

Property: There exists a bound *m* such that every variable introduced is *m*-consumed.

<pre>let kalman = stream { init = 0.0; step (pre_x, obs) = { let x = sample(gaussian(pre_x,1)) in let y = sample(gaussian(x,1)) in observe (y, obs); (x, x) } </pre>	$x0 \leftarrow nil ::$ $y0 \leftarrow x0::$ $y0 \leftarrow x0::$ y0 is 1-consumed $x1 \leftarrow x0::$ $y1 \leftarrow x1::$ observe y1:: $x2 \leftarrow x1::$ $y2 \leftarrow y1::$ observe y2::
} }	•••

The Unseparated Paths Property

Property: there exists an *n* such that at any time step t, no variable in the program state at t starts an unseparated path longer than *n*

```
let kalman = stream {
  init = 0.0;
 step (pre_x, obs) = {
   let x = sample(gaussian(pre_x,1)) in
   let y = sample(gaussian(x,1)) in
   observe (y, obs);
   (x, x)
```



← x0::

↔ x0 ::

 \leftarrow x1 ::

← x1 ::

← y1 ።

path longer than 2

- Can approximate with a larger upper bound
- Analysis passes when longest path converges

<pre>let kalman = stream {</pre>	$x0 \leftarrow nil ::$	(x
init = 0.0;	y0 ← x0∷	(x
<pre>step (pre_x, obs) = {</pre>	observe y0 ::	(x
<pre>let x = sample(gaussian(pre_x,1)) in</pre>	x1 ← x0 ::	(x
<pre>let y = sample(gaussian(x,1)) in</pre>	y1 ← x1 ∷	(x
observe (y, obs);	observe y1 ::	(x
(x, x)	$x2 \leftarrow x1 ::$	(x
}	y2 ← y1 ∷	(x
}	observe y2 ::	(x

Program State

				<i>m</i> -consumed		unsep.	paths	
:= 0	x0			output	actual	output	actual	(
		Analysis is	Kalman	\checkmark	\checkmark	\checkmark	\checkmark	
: = 1		nrecise \ Kalman Hold-First	\checkmark	\checkmark	X	X		
	x1	precise	Gaussian Random Walk	X	X	\checkmark	\checkmark	
; = 2			Robot	\checkmark	\checkmark	\checkmark	\checkmark	
	v2	x2 Memory is Probabilistically Bounded Outlier	Coin	\checkmark	\checkmark	\checkmark	\checkmark	
	Λ Ζ		Gaussian-Gaussian	\checkmark	\checkmark	\checkmark	\checkmark	
			Outlier	X	X	\checkmark	\checkmark	
			MTT	X	×	\checkmark	\checkmark	
			SLAM	X	\checkmark	\checkmark	\checkmark	
		Memory is						
		Always /						
		Bounded						

Results



y0 ← x0::



	Longest
	Unsep. Path
	(x0, x0), 0
	(x0, y0), 1
•	(x0, y0), 1
	(x0, x1), 1
	(x1, y1), 1
•	(x1, y1), 1
	(x1, x2), 1
	(x2, y2), 1
•	(x2, y2), 1

